

1. The *ACTION* function takes as arguments a state i and a terminal a ($\$,$ the input endmarker). The value of $ACTION[i,a]$ can have one of four forms:
 - (a) Shift j where j is a state. The action taken by the parser effectively shifts input a to the stack, but uses state j to represent a .
 - (b) Reduce $A \rightarrow \beta$. The action of the parser effectively reduces β on the top of the stack to head A .
 - (c) Accept. The parser accepts the input and finishes parsing.
 - (d) Error. The parser discovers an error in its input and takes some corrective action.
2. We extend the *GOTO* function, defined on sets of items, to states: if $GOTO[I_i, A] = I_j$, then *GOTO* also maps a state i and a nonterminal A to state j .

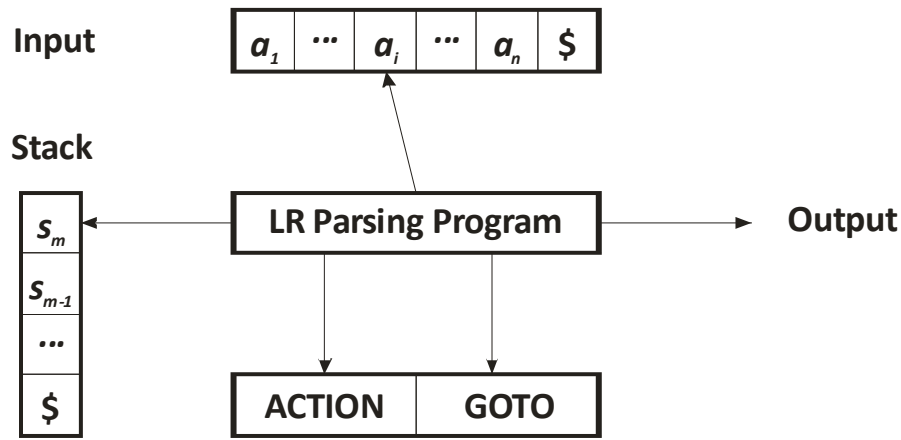


Figure 1. LR Parser Model

	<i>left side</i>		<i>right side</i>
1	E	$\rightarrow$	$E + T$
2	E	$\rightarrow$	T
3	T	$\rightarrow$	$T * F$
4	T	$\rightarrow$	F
5	F	$\rightarrow$	(E)
6	F	$\rightarrow$	id

Table 1. Set of productions expressions

let a be the first symbol of $w\$$

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while (1) {
    let  $s$  be the state on top of the stack;
    if ( $ACTION[s,a]==shift\ t$ ) {
        push  $t$  onto the stack
        let  $a$  be the next input symbol;
    } else if ( $ACTION[s,a]==reduce\ A \rightarrow \beta$ ) {
        pop  $|\beta|$  symbols off the stack;
        let state  $t$  now be on top of the stack;
        push  $GOTO[t,A]$  onto the stack
        output the production  $A \rightarrow \beta$  ;
    } else if ( $ACTION[s,a]==accept$ ) break; //parsing is done
    else error();
}

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	STACK	SYMBOLS	INPUT	ALGORITHM	ACTION
(1)	0		id * id + id \$	$ACTION[0, id]=s5$	shift
(2)	0 5	id	* id + id \$	$ACTION[5, *]=r6$ $GOTO[0, F]=3$	reduce by $F \rightarrow id$
(3)	0 3	F	* id + id \$	$ACTION[3, *]=r4$ $GOTO[0, T]=2$	reduce by $T \rightarrow F$
(4)	0 2	T	* id + id \$	$ACTION[2, *]=s7$	shift
(5)	0 2 7	$T *$	id + id \$	$ACTION[7, id]=s5$	shift
(6)	0 2 7 5	$T * id$	+ id \$	$ACTION[5, +]=r6$ $GOTO[7, F]=10$	reduce by $F \rightarrow id$
(7)	0 2 7 10	$T * F$	+ id \$	$ACTION[10, +]=r3$ $GOTO[0, T]=2$	reduce by $T \rightarrow T * F$
(8)	0 2	T	+ id \$	$ACTION[2, +]=r2$ $GOTO[0, E]=1$	reduce by $E \rightarrow T$
(9)	0 1	E	+ id \$	$ACTION[1, +]=s6$	shift
(10)	0 1 6	$E +$	id \$	$ACTION[6, id]=s5$	shift
(11)	0 1 6 5	$E + id$	\$	$ACTION[5, \$]=r6$ $GOTO[6, F]=3$	reduce by $F \rightarrow id$
(12)	0 1 6 3	$E + F$	\$	$ACTION[3, \$]=r4$ $GOTO[6, T]=9$	reduce by $T \rightarrow F$
(13)	0 1 6 9	$E + T$	\$	$ACTION[9, \$]=r1$ $GOTO[0, E]=1$	reduce by $E \rightarrow E + T$
(14)	0 1	E	\$	$ACTION[1, \$]=acc$	accept

STATE	ACTION						GOTO		
	id	+	*	(	)	\$	<i>E</i>	<i>T</i>	<i>F</i>
0	s5			s4			1	2	3
1		s6				acc			
2		r2	s7		r2	r2			
3		r4	r4		r4	r4			
4	s5			s4			8	2	3
5		r6	r6		r6	r6			
6	s5			s4				9	3
7	s5			s4					10
8		s6			s11				
9		r1	s7		r1	r1			
10		r3	r3		r3	r3			
11		r5	r5		r5	r5			

STATE	ACTION						GOTO		
	id	+	*	(	)	\$	E	T	F
0	s5			s4			1	2	3
1		s6				acc			
2		r2	s7		r2	r2			
3		r4	r4		r4	r4			
4	s5			s4			8	2	3
5		r6	r6		r6	r6			
6	s5			s4				9	3
7	s5			s4					10
8		s6			s11				
9		r1	s7		r1	r1			
10		r3	r3		r3	r3			
11		r5	r5		r5	r5			

	STACK	SYMBOLS	INPUT	ALGORITHM	ACTION
(1)	0		id + id * id \$	ACTION[0,id]=s5	shift
(2)	0 5	id	+ id * id \$	ACTION[5,+]=r6 GOTO[0,F]=3	reduce by $F \rightarrow \text{id}$
(3)	0 3	F	+ id * id \$	ACTION[3,+]=r4 GOTO[0,T]=2	reduce by $T \rightarrow F$
(4)	0 2	T	+ id * id \$	ACTION[2,+]=r2 GOTO[0,E]=1	reduce by $E \rightarrow T$
(5)	0 1	E	+ id * id \$	ACTION[1,+]=s6	shift
(6)	0 1 6	E +	id * id \$	ACTION[6,id]=s5	shift
(7)	0 1 6 5	E + id	* id \$	ACTION[5,*]=r6 GOTO[6,F]=3	reduce by $F \rightarrow \text{id}$
(8)	0 1 6 3	E + F	* id \$	ACTION[3,*]=r4 GOTO[6,T]=9	reduce by $T \rightarrow F$
(9)	0 1 6 9	E + T	* id \$	ACTION[9,*]=s7	shift
(10)	0 1 6 9 7	E + T *	id \$	ACTION[7,id]=s5	shift
(11)	0 1 6 9 7 5	E + T * id	\$	ACTION[5,\$]=r6 GOTO[7,F]=10	reduce by $F \rightarrow \text{id}$
(12)	0 1 6 9 7 10	E + T * F	\$	ACTION[10,\$]=r3 GOTO[6,T]=9	reduce by $T \rightarrow T * F$
(13)	0 1 6 9	E + T	\$	ACTION[9,\$]=r1 GOTO[0,E]=1	reduce by $E \rightarrow E + T$
(14)	0 1	E	\$	ACTION[1,\$]=acc	accept

$(s_0 s_1 \cdots s_m, a_i a_{i+1} \cdots a_n \$)$

$X_1 X_2 \cdots X_m a_i a_{i+1} \cdots a_n$